



**FINITE AUTOMATA**

ประเด็นที่ควรสนใจ  
คำถาม

How to know which word is in the given language ?  
How to know if any given word is in the given language ?

## FINITE AUTOMATA

# ออโตมาตาคำจำกัด นิยาม

A finite automaton or finite state machine (abbreviated FA) is a 5-tuple  $(Q, \Sigma, q_0, \delta, A)$  where

- $Q$  means a finite set of states.
- $\Sigma$  is a finite input alphabet.
- $q_0 \in Q$  named Initial state.
- $A \subseteq Q$ ,  $A$  is the set of all accepted states.
- $\delta$  is a function from  $Q \times \Sigma$  to  $Q$ , called transition function.

## FINITE AUTOMATA

# ออโตมาตาคำจำกัดตัวอย่าง

Suppose that  $\Sigma = \{0, 1\}$ , and  $Q = \{q_0, q_1, q_2\}$ .

The transition function  $\delta$  is defined as follow;

$$\delta(q_0, 0) = q_0$$

$$\delta(q_0, 1) = q_1$$

$$\delta(q_1, 0) = q_2$$

$$\delta(q_2, 0) = q_0$$

$$\delta(q_2, 1) = q_1$$

The initial state is  $q_0$ .

The accepted or final state is  $q_2$ .

Example:

Consider this string 01001.

Start at the initial state  $q_0$ ,

$\delta(q_0, 0)$  gives state  $q_0$ .

$\delta(q_0, 1)$  gives state  $q_1$ .

$\delta(q_1, 0)$  gives state  $q_2$ .

$\delta(q_2, 0)$  gives state  $q_0$ .

$\delta(q_0, 1)$  gives state  $q_1$ .

It ends at the state  $q_1$  which is not the accepted state.

We say that it **rejects** the input, otherwise we say it accepts.

FINITE AUTOMATA

# ออโตมาตารูปตัวอย่าง

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$$\delta(q_0, 1) = q_1$$

$$\delta(q_1, 0) = q_0$$

$$\delta(q_1, 1) = q_1$$

The initial state is  $q_0$ .

The accepted state is  $q_1$ .

Example:

Consider this string 0101.

Start:  $\delta(q_0, 0) \rightarrow q_0$ .

$\delta(q_0, 1) \rightarrow q_1$ .

$\delta(q_1, 0) \rightarrow q_0$ .

$\delta(q_0, 1) \rightarrow q_1$ .

Then it **accepts** 0101.

Consider 0110.

Start:  $\delta(q_0, 0) \rightarrow q_0$ .

$\delta(q_0, 1) \rightarrow q_1$ .

$\delta(q_1, 1) \rightarrow q_1$ .

$\delta(q_1, 0) \rightarrow q_0$ .

Then it **rejects** 0110.

FINITE AUTOMATA

# ออโตมาตารูปตัวอย่าง

Suppose that  $\Sigma = \{ 0, 1 \}$ , and  $Q = \{ q_0, q_1 \}$ .

The transition function  $\delta$  is defined as follow;

$$\delta(q_0, 0) = q_0$$

$$\delta(q_0, 1) = q_1$$

$$\delta(q_1, 0) = q_0$$

$$\delta(q_1, 1) = q_1$$

The initial state is  $q_0$ .

The accepted state is  $q_1$ .

TABLE:

State/input	0	1
$q_0$	$q_0$	$q_1$
$q_1$	$q_0$	$q_1$

REPRESENTATION BY TABLE

FINITE AUTOMATA

# ออโตมาตาคำจำกัดตัวอย่าง

Suppose that  $\Sigma = \{ 0, 1 \}$ , and  $Q = \{ q_0, q_1 \}$ .  
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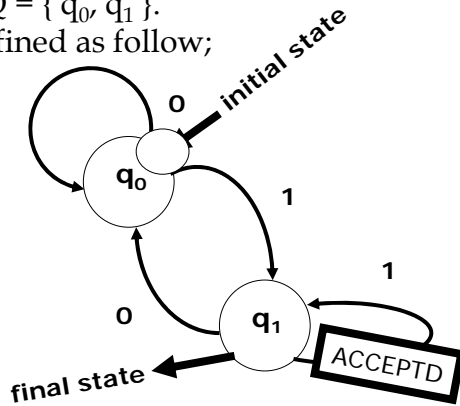
$$\delta(q_1, 1) = q_1$$

The initial state is  $q_0$ .

The accepted state is  $q_1$ .

**Input string is 0101**

REPRESENTATION BY DIAGRAM



FINITE AUTOMATA

# ออโตมาตาคำจำกัด นิยาม

accepted word  $\neq$  rejected word

The set of all strings that leave the finite state machine in the final state is called the language defined by the finite state machine or the language associated with this machine.

language-recognizer

FINITE AUTOMATA

# ออโตมาตาคำจำกัด นิยาม

ACCEPTANCE BY FA

accepted word  $\neq$  rejected word

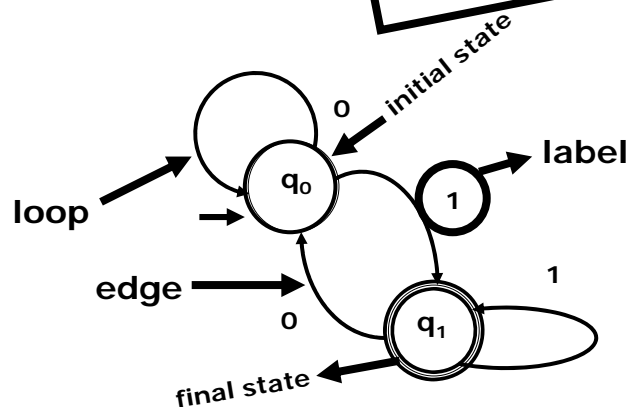
The set of all strings that leave the finite state machine in the final state is called the language defined by the finite state machine or the language associated with this machine.

language-recognizer

$$L(M) = \{ x \in \Sigma^* \mid x \text{ is accepted by } M \}$$

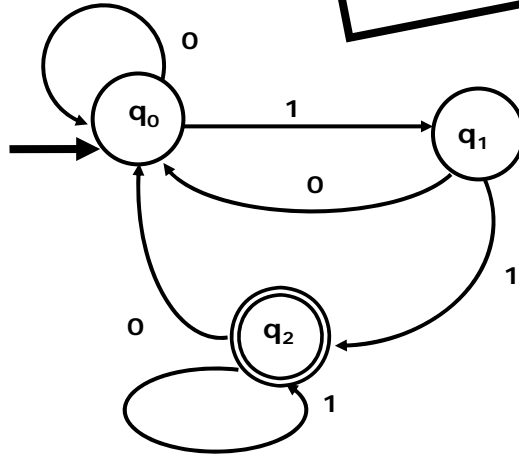
FINITE AUTOMATA

# ออโตมาตาคำจำกัด นิยาม



FINITE AUTOMATA

# ออโตมาตาจำกัด ทดลอง

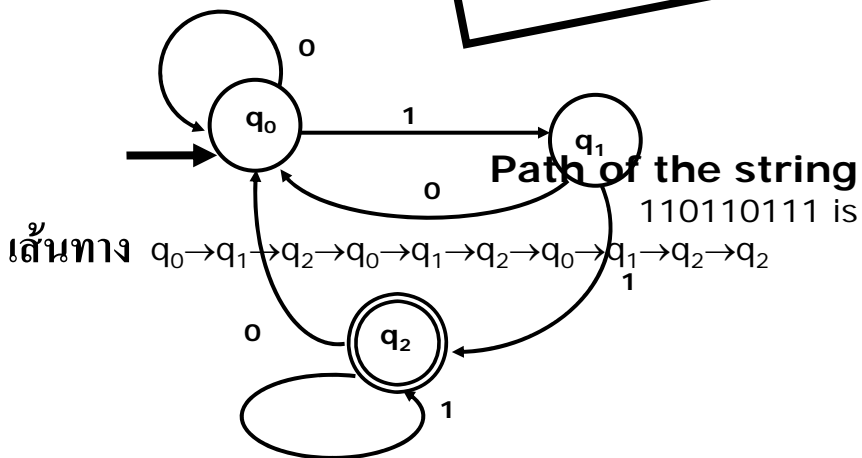


GIVEN

- ~~111001010~~
- 110110111
- 000000111
- ~~011011100~~
- ~~010010101~~

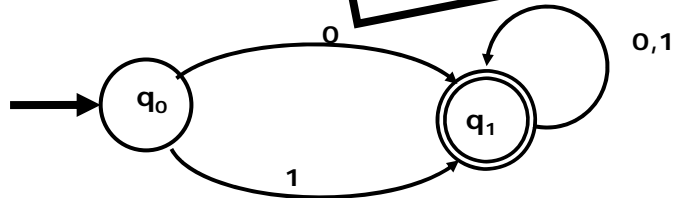
FINITE AUTOMATA

# ออโตมาตาจำกัด นิยาม



FINITE AUTOMATA

# ออโตมาตารูปตัวอย่าง



Find the language accepted by this finite machine.

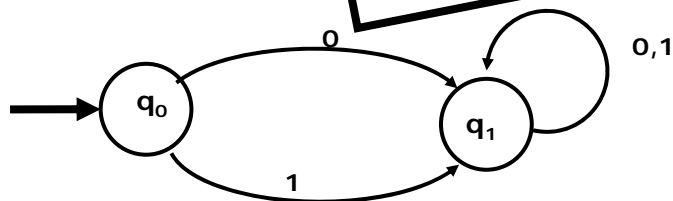
$(0+1)(0+1)^*$

or

$(0+1)^+$

FINITE AUTOMATA

# ออโตมาตารูปตัวอย่าง

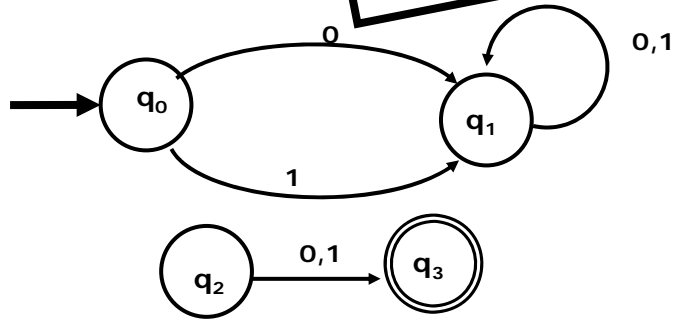


There is no final state.

AUTOMATA WITH NO LANGUAGE

FINITE AUTOMATA

# ออโตมาตารำกั้ดตัวอย่าง

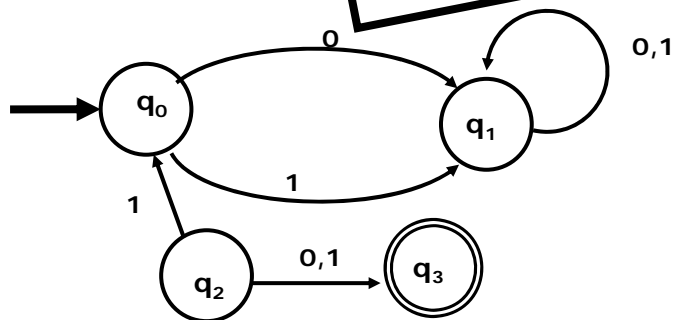


The graph is disconnected.

AUTOMATA WITH NO LANGUAGE

FINITE AUTOMATA

# ออโตมาตารำกั้ดตัวอย่าง



There is no path from the initial state to the final state.

AUTOMATA WITH NO LANGUAGE



FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Studies finite automata for two different angles:

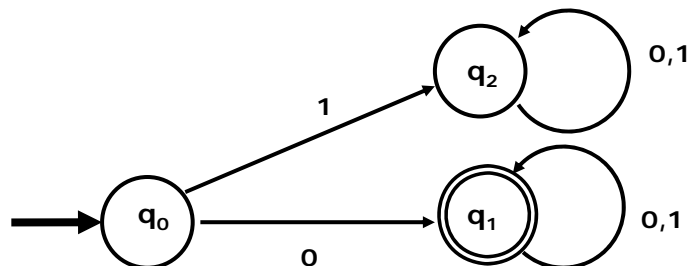
- Given a language, can we build a machine for it ?
- Given a machine, can we deduce its language ?

MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Suppose we want to build a finite machine that accepts all words in the language  $0(1+0)^*$ .

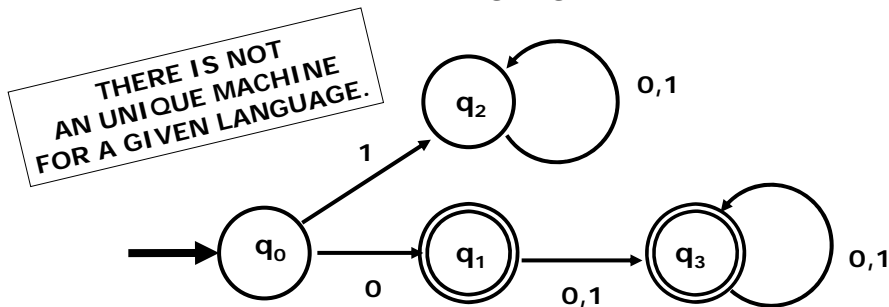


MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Suppose we want to build a finite machine that accepts all words in the language  $0(1+0)^*$ .



MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

From example, we can ask a question

Is there always at least one finite automaton that accepts each possible language ?

This is related to the question: Can all languages be recognized by a finite automaton ?

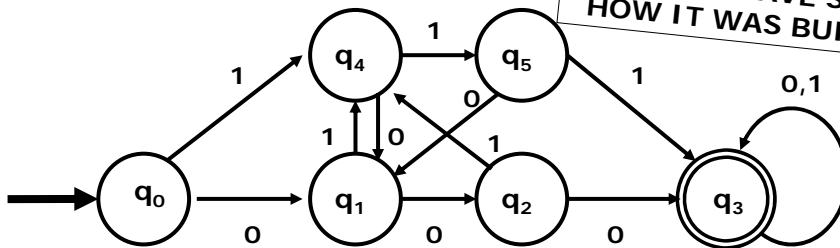
MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Suppose we want to build a finite machine that accepts all words in the language containing a triple characters, (i.e., 000 or 111).

EASY ! WE HAVE SEEN HOW IT WAS BUILT

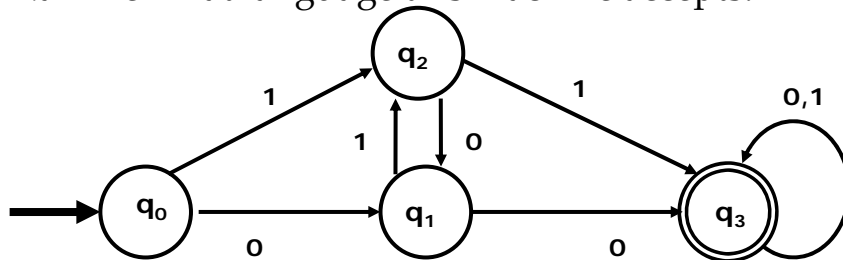


MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Examine what language this machine accepts.



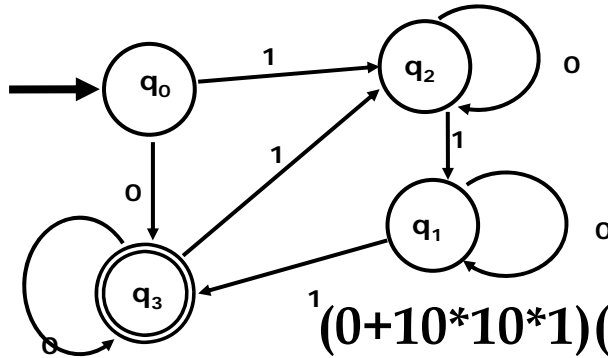
$$(0+1)^*(00+11)(0+1)^*$$

MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Examine what language this machine accepts.



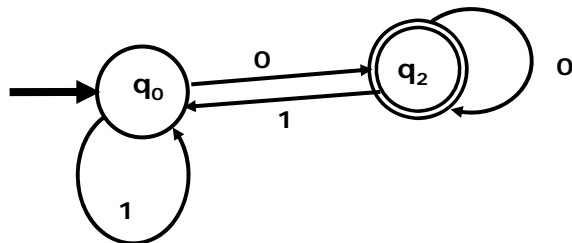
$$1(0+10^*10^*1)(0+10^*10^*1)^*$$

MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Examine what language this machine accepts.



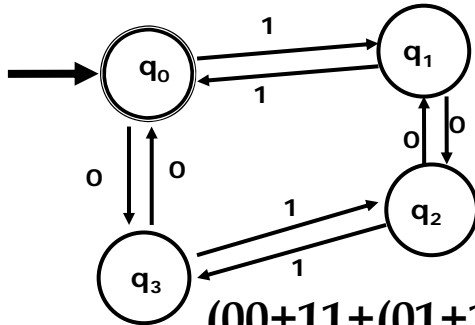
$$(0+1)^*0$$

MACHINES AND THEIR LANGUAGES

FINITE AUTOMATA

# ออโตมาตาคำจำกัด ภาษา

Examine what language this machine accepts.



**EVEN-EVEN  
LANGUAGE**

$(00+11+(01+10)(00+11)^*(01+10))^*$

MACHINES AND THEIR LANGUAGES