



Graphs & Trees (Lecture 2)

Today's Topics

- Graph Representation
- Graph Isomorphism
- Graph Connectivity / Paths



Representing Graphs

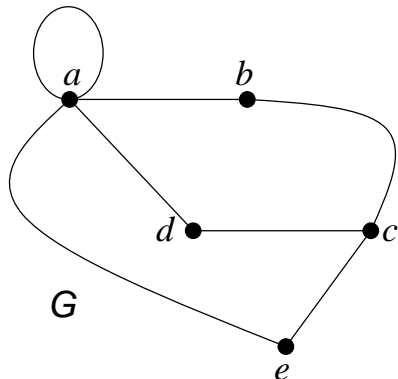
We can use:

- Adjacency Lists
- Adjacency Matrices
- Incidence Matrices



Adjacency Lists

- represent graphs with no multiple edges
- specify vertices that are adjacent to each vertex

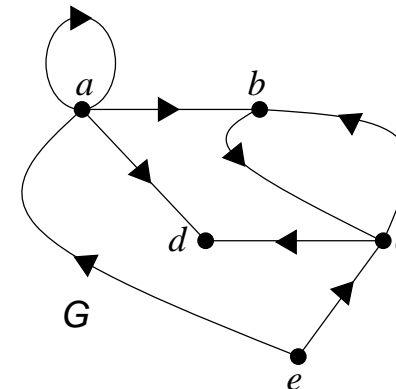


vertex	adjacent vertices
<i>a</i>	<i>a b d e</i>
<i>b</i>	<i>a c</i>
<i>c</i>	<i>b d e</i>
<i>d</i>	<i>a c</i>
<i>e</i>	<i>a c</i>



Adjacency Lists

- For directed graph, specify terminal vertices of each vertex



initial vertex	terminal vertices
<i>a</i>	<i>a b d</i>
<i>b</i>	<i>c</i>
<i>c</i>	<i>b d</i>
<i>d</i>	-
<i>e</i>	<i>a c</i>



Adjacency Matrices

$$G = (V, E), V = \{v_1, v_2, \dots, v_n\}$$

$$A = [a_{ij}]$$

Directed Graph

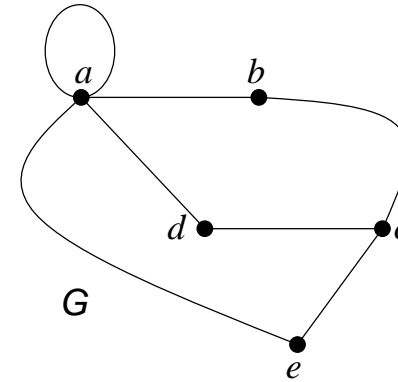
a_{ij} = Number of edges corresponding to $\{v_i, v_j\}$

Undirected Graph

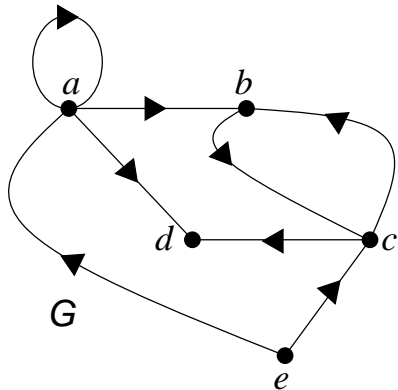
a_{ij} = Number of edges corresponding to (v_i, v_j)



Adjacency Matrices



Adjacency Matrices



Incidence Matrices

$$G = (V, E), V = \{v_1, v_2, \dots, v_n\}, E = \{e_1, e_2, \dots, e_m\}$$

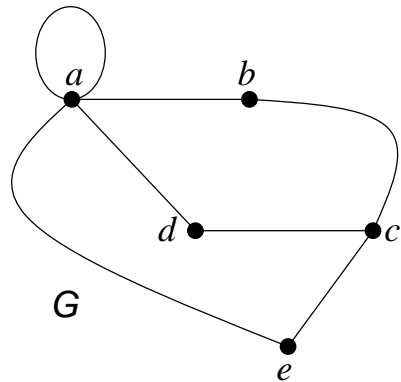
Undirected Graph

$$M = [m_{ij}]$$

$$m_{ij} = \begin{cases} 1 & \text{when } e_j \text{ is incident with } v_i \\ 0 & \text{otherwise.} \end{cases}$$



Incidence Matrices



Isomorphism of Graphs

$G_1=(V_1,E_1)$ and $G_2=(V_2,E_2)$ are two simple graphs.
 G_1 and G_2 are **isomorphic** when:

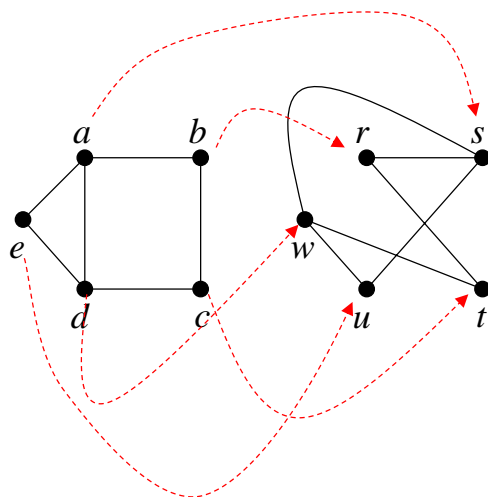
There is a 1-to-1 and onto function f from V_1 to V_2 with the property that

$$\forall a, b \in V_1 \quad a \text{ and } b \text{ are adjacent in } G_1 \leftrightarrow f(a) \text{ and } f(b) \text{ are adjacent in } G_2.$$

f is called an **isomorphism**.



Isomorphism of Graphs



Isomorphism, f
 $f(a) = s, f(b) = r$
 $f(c) = t, f(d) = w$
 $f(e) = u$

$\{a,b\} \rightarrow \{r,s\}$
 $\{b,c\} \rightarrow \{r,t\}$
 $\{c,d\} \rightarrow \{t,w\}$
 $\{d,a\} \rightarrow \{u,w\}$
 $\{d,e\} \rightarrow \{s,u\}$
 $\{a,e\} \rightarrow \{s,w\}$



Isomorphism of Graphs

- To show that two graphs are isomorphic



Find an isomorphism.

- To show that two graphs are NOT isomorphic



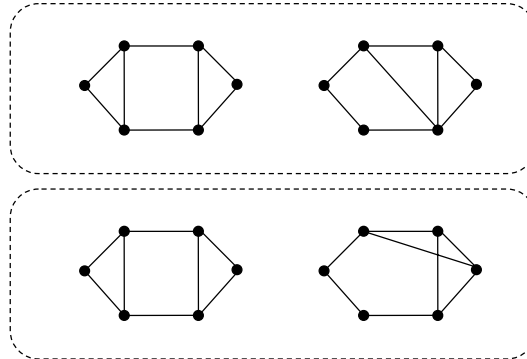
Argue that no isomorphism exists.



Graph Invariants

- A property preserved by isomorphism of graphs is called a **graph invariant**.

- E.g.:
- ♥ Number of vertices
 - ♥ Number of edges
 - ♥ Degrees of vertices
 - ♥ Adjacency between vertices with specified degrees

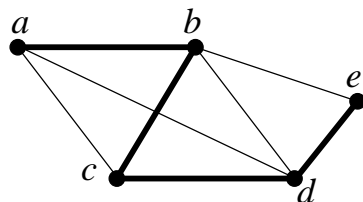


Paths

- A **path** is a sequence of edges that begins with a vertex of a graph and travels from vertex to vertex along edges of the graph.
- The **length** of a path is the number of edges in that path.
- A path is a **circuit** if it begins and ends at the same vertex and the length is not 0.
- A path is **simple** if it does not contain the same edge more than once.



Paths



$(\{a,b\}, \{b,c\}, \{c,d\}, \{d,e\})$ is a simple path of length 4.

When there are no multiple paths, we can represent a path by the sequence of vertices it pass through.

$$(\{a,b\}, \{b,c\}, \{c,d\}, \{d,e\}) \Rightarrow (a,b,c,d,e)$$

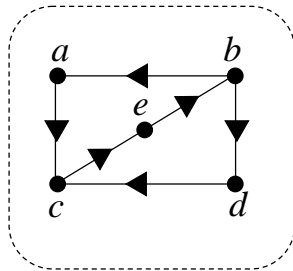
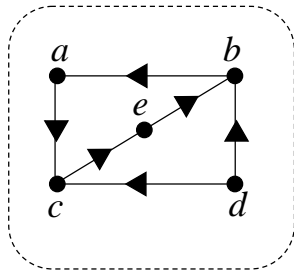
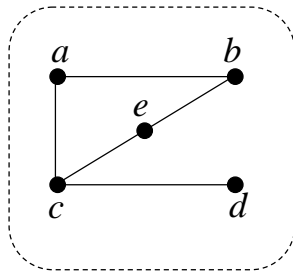
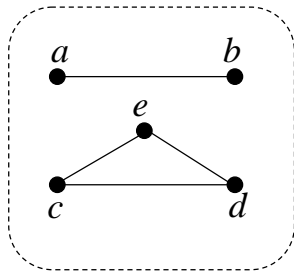


Connectedness

An undirected graph is called **connected** \leftrightarrow There is a path between every pair of distinct vertices of the graph.

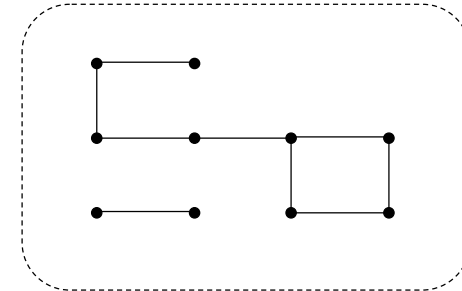
An directed graph is **strongly connected** \leftrightarrow There is a path from a to b and from b to a whenever a and b are distinct vertices in the graph.

An directed graph is **weakly connected** \leftrightarrow There is a path between every pair of distinct vertices of the graph. (Disregard the edge directions.)



Connected Components

- A **connected component** of a graph is maximal connected subgraph of that graph.

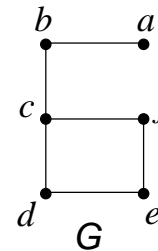


Cut Vertex / Cut Edge

- A vertex is called a **cut vertex** (or **articulation point**) if the removal of that vertex along with all edges incident with it produces a subgraph with more connected components.
- An edge whose removal produces a subgraph with more connected components is called a **cut edge** (or **bridge**).



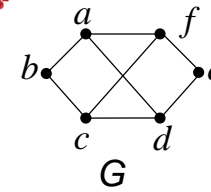
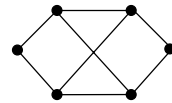
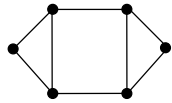
Find all cut vertices and cut edges in G .





Paths and Isomorphism

- ♥ The existence of a simple circuit of a particular length is a graph invariant.
- ♥ Paths are sometimes useful in finding an isomorphism.

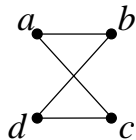


Show that a simple circuit of length 3 does not exist in G .



Counting Paths Between Vertices

How many paths of length 4 are there from a to d in the following graph?



Counting Paths Between Vertices

The number of paths of a certain length between two vertices in a graph can also be determined from the graph's adjacency matrix.

Find out by yourself 😊

