

# **Chapter #2: Two-Level Combinational Logic**

***Contemporary Logic Design***

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## Logic Functions: Boolean Algebra

Algebraic structure consisting of:

set of elements  $B$

binary operations  $\{+, \cdot\}$

unary operation  $\{\prime\}$

such that the following axioms hold:

1.  $B$  contains at least two elements,  $a, b$ , such that  $a \neq b$
2. **Closure**  $a, b$  in  $B$ ,
  - (i)  $a + b$  in  $B$
  - (ii)  $a \cdot b$  in  $B$
3. **Commutative Laws:**  $a, b$  in  $B$ ,
  - (i)  $a + b = b + a$
  - (ii)  $a \cdot b = b \cdot a$
4. **Identities:**  $0, 1$  in  $B$ 
  - (i)  $a + 0 = a$
  - (ii)  $a \cdot 1 = a$
5. **Distributive Laws:**
  - (i)  $a + (b \cdot c) = (a + b) \cdot (a + c)$
  - (ii)  $a \cdot (b + c) = a \cdot b + a \cdot c$
6. **Complement:**
  - (i)  $a + a' = 1$
  - (ii)  $a \cdot a' = 0$

# Logic Functions: Boolean Algebra

$B = \{0,1\}$ ,  $+$  = OR,  $\cdot$  = AND,  $'$  = NOT is a Boolean Algebra

must verify that the axioms hold:

E.g., Commutative Law:

$$0 + 1 = 1 + 0?$$

$$1 = 1$$

$$0 \cdot 1 = 1 \cdot 0?$$

$$0 = 0$$

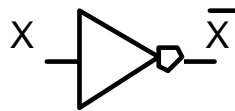
**Theorem:** any Boolean function that can be expressed as a truth table can be written as an expression in Boolean Algebra using  $'$ ,  $+$ ,  $\cdot$

Review  
from  
Chapter 1

Description

If  $X = 0$  then  $X' = 1$   
If  $X = 1$  then  $X' = 0$

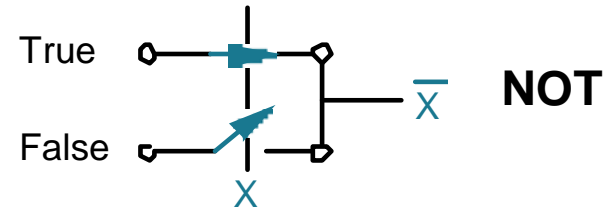
Gates



Truth Table

X	$\bar{X}$
0	1
1	0

Switches



Description

$Z = 1$  if  $X$  and  $Y$   
are both 1

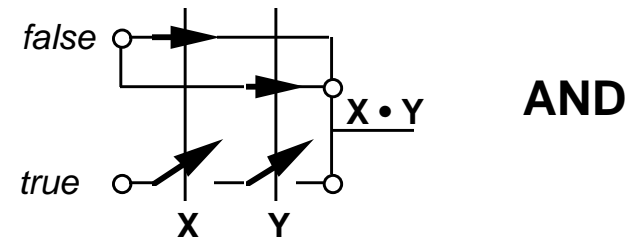
Gates



Truth Table

X	Y	Z
0	0	0
0	1	0
1	0	0
1	1	1

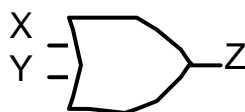
Switches



Description

$Z = 1$  if  $X$  or  $Y$   
(or both) are 1

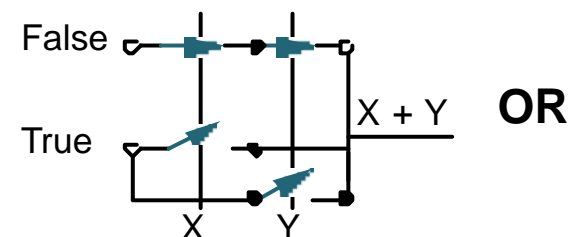
Gates



Truth Table

X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	1

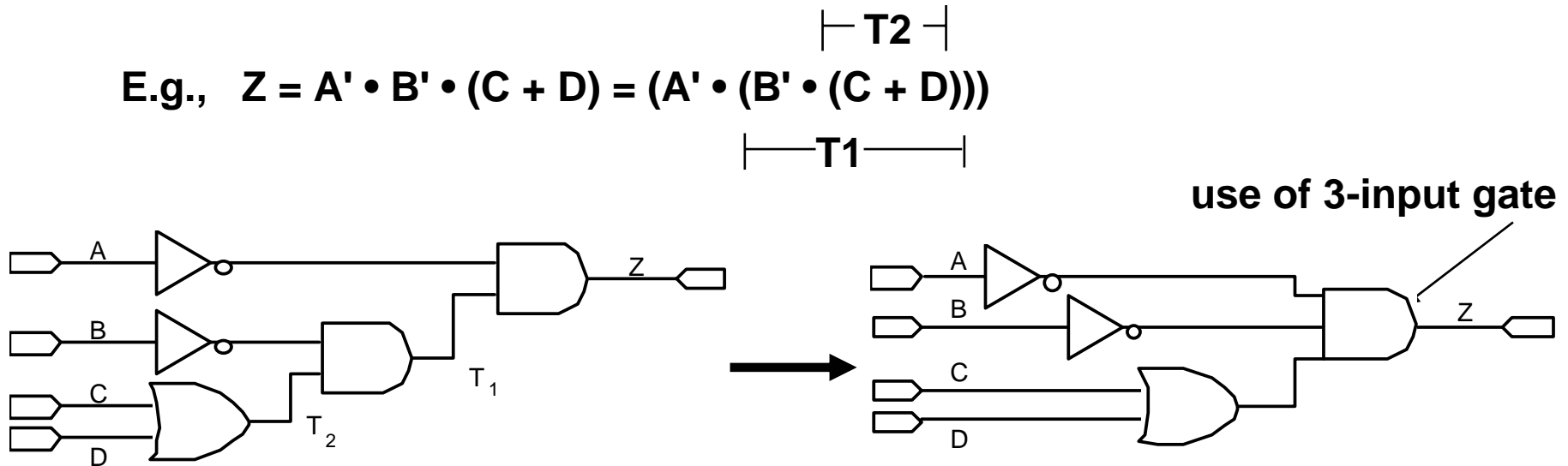
Switches



## Logic Functions: From Expressions to Gates

More than one way to map an expression to gates

E.g.,  $Z = A' \cdot B' \cdot (C + D) = (A' \cdot (B' \cdot (C + D)))$



**Literal:** each appearance of a variable or its complement in an expression

E.g.,  $Z = A B' C + A' B + A' B C' + B' C$

3 variables, 10 literals

# Logic Functions: NAND, NOR, XOR, XNOR

## 16 functions of two variables:

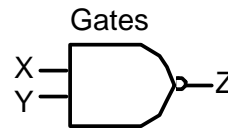
X	Y	F0	F1	F2	F3	F4	F5	F6	F7	F8	F9	F10	F11	F12	F13	F14	F15
0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1	0	0	0	1	1	1	1	1
1	0	0	0	1	1	0	0	1	1	0	0	1	0	0	0	1	1
1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1

0
 $X \cdot Y$ 
 $X$ 
 $Y$ 
 $X + Y$ 
 $\bar{Y}$ 
 $\bar{X}$ 
1

**X, X', Y, Y', X•Y, X+Y, 0, 1 only  
half of the possible functions**

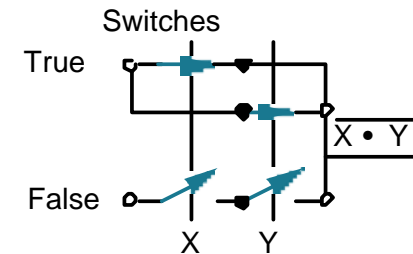
### NAND

Description  
 $Z = 1$  if X is 0  
or Y is 0



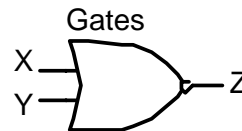
Truth T able

X	Y	Z
0	0	1
0	1	1
1	0	1
1	1	0



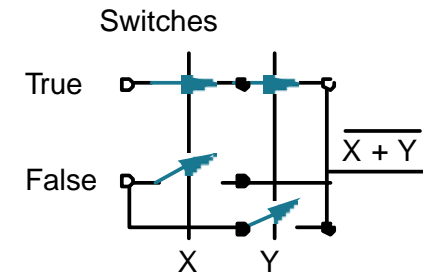
### NOR

Description  
 $Z = 1$  if both X  
and Y are 0



Truth T able

X	Y	Z
0	0	1
0	1	0
1	0	0
1	1	0



## Logic Functions: NAND, NOR Implementation

NAND, NOR gates far outnumber AND, OR in typical designs  
*easier to construct in the underlying transistor technologies*

Any Boolean expression can be implemented by NAND, NOR, NOT gates

In fact, NOT is superfluous  
(NOT = NAND or NOR with both inputs tied together)

X	Y	X NOR Y
0	0	1
1	1	0

X	Y	X NAND Y
0	0	1
1	1	0

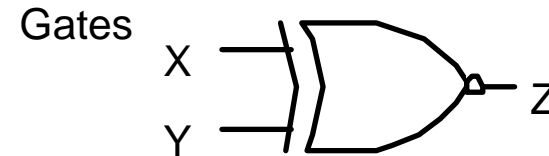
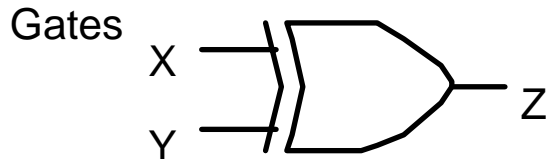
## Logic Functions: XOR, XNOR

**XOR: X or Y but not both ("inequality", "difference")**

**XNOR: X and Y are the same ("equality", "coincidence")**

Description  
Z = 1 if X has a different  
value than Y

Description  
Z = 1 if X has the same  
value as Y



Truth T able

X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	0

(a) XOR

$$X \oplus Y = X Y' + X' Y$$

Truth T able

X	Y	Z
0	0	1
0	1	0
1	0	0
1	1	1

(b) XNOR

$$\overline{X \oplus Y} = X Y + X' Y'$$

## Gate Logic: Laws of Boolean Algebra

**Duality:** a dual of a Boolean expression is derived by replacing AND operations by ORs, OR operations by ANDs, constant 0s by 1s, and 1s by 0s (literals are left unchanged).

***Any statement that is true for an expression is also true for its dual!***

### Useful Laws/Theorems of Boolean Algebra:

#### ***Operations with 0 and 1:***

1.  $X + 0 = X$

2.  $X + 1 = 1$

1D.  $X \cdot 1 = X$

2D.  $X \cdot 0 = 0$

#### ***Idempotent Law:***

3.  $X + X = X$

3D.  $X \cdot X = X$

#### ***Involution Law:***

4.  $(X')' = X$

#### ***Laws of Complementarity:***

5.  $X + X' = 1$

5D.  $X \cdot X' = 0$

#### ***Commutative Law:***

6.  $X + Y = Y + X$

6D.  $X \cdot Y = Y \cdot X$



## Gate Logic: Laws of Boolean Algebra (cont)

### *Associative Laws:*

$$7. (X + Y) + Z = X + (Y + Z) \\ = X + Y + Z$$

$$7D. (X \cdot Y) \cdot Z = X \cdot (Y \cdot Z) \\ = X \cdot Y \cdot Z$$

### *Distributive Laws:*

$$8. X \cdot (Y + Z) = (X \cdot Y) + (X \cdot Z)$$

$$8D. X + (Y \cdot Z) = (X + Y) \cdot (X + Z)$$

### *Simplification Theorems:*

$$9. X \cdot Y + X \cdot Y' = X$$

$$9D. (X + Y) \cdot (X + Y') = X$$

$$10. X + X \cdot Y = X$$

$$10D. X \cdot (X + Y) = X$$

$$11. (X + Y') \cdot Y = X \cdot Y$$

$$11D. (X \cdot Y') + Y = X + Y$$

### *DeMorgan's Law:*

$$12. (X + Y + Z + \dots)' = X' \cdot Y' \cdot Z' \cdot \dots \quad 12D. (X \cdot Y \cdot Z \cdot \dots)' = X' + Y' + Z' + \dots$$

$$13. \{F(X_1, X_2, \dots, X_n, 0, 1, +, \cdot)\}' = \{F(X_1', X_2', \dots, X_n', 1, 0, \cdot, +)\}$$

### *Duality:*

$$14. (X + Y + Z + \dots)^D = X \cdot Y \cdot Z \cdot \dots \quad 14D. (X \cdot Y \cdot Z \cdot \dots)^D = X + Y + Z + \dots$$

$$15. \{F(X_1, X_2, \dots, X_n, 0, 1, +, \cdot)\}^D = \{F(X_1, X_2, \dots, X_n, 1, 0, \cdot, +)\}$$

### *Theorems for Multiplying and Factoring:*

$$16. (X + Y) \cdot (X' + Z) = X \cdot Z + X' \cdot Y \quad 16D. X \cdot Y + X' \cdot Z = (X + Z) \cdot (X' + Y)$$

### *Consensus Theorem:*

$$17. (X \cdot Y) + (Y \cdot Z) + (X' \cdot Z) = \\ X \cdot Y + X' \cdot Z$$

$$17D. (X + Y) \cdot (Y + Z) \cdot (X' + Z) = \\ (X + Y) \cdot (X' + Z)$$

## Gate Logic: Laws of Boolean Algebra

*Proving theorems via axioms of Boolean Algebra:*

E.g., prove the theorem:  $X \cdot Y + X \cdot Y' = X$

E.g., prove the theorem:  $X + X \cdot Y = X$

**Gate Logic: Laws of Boolean Algebra**

***Proving theorems via axioms of Boolean Algebra:***

E.g., prove the theorem:  $X \cdot Y + X \cdot Y' = X$

*distributive law (8)*       $X \cdot Y + X \cdot Y' = X \cdot (Y + Y')$

*complementary law (5)*       $X \cdot (Y + Y') = X \cdot (1)$

*identity (1D)*       $X \cdot (1) = X$

E.g., prove the theorem:  $X + X \cdot Y = X$

*identity (1D)*       $X + X \cdot Y = X \cdot 1 + X \cdot Y$

*distributive law (8)*       $X \cdot 1 + X \cdot Y = X \cdot (1 + Y)$

*identity (2)*       $X \cdot (1 + Y) = X \cdot (1)$

*identity (1)*       $X \cdot (1) = X$

## Gate Logic: Laws of Boolean Algebra

### DeMorgan's Law

$$(X + Y)' = X' \cdot Y'$$

**NOR is equivalent to AND  
with inputs complemented**

X	Y	$\bar{X}$	$\bar{Y}$	$\overline{X+Y}$	$\overline{X \cdot Y}$
0	0	1	1	1	1
0	1	1	0	0	0
1	0	0	1	0	0
1	1	0	0	0	0

$$(X \cdot Y)' = X' + Y'$$

**NAND is equivalent to OR  
with inputs complemented**

X	Y	$\bar{X}$	$\bar{Y}$	$\overline{X \cdot Y}$	$\overline{X+Y}$
0	0	1	1	1	1
0	1	1	0	1	1
1	0	0	1	1	1
1	1	0	0	0	0

**DeMorgan's Law can be used to convert AND/OR expressions  
to OR/AND expressions**

**Example:**

$$Z = A' B' C + A' B C + A B' C + A B C'$$

$$Z' = (A + B + C') \cdot (A + B' + C') \cdot (A' + B + C') \cdot (A' + B' + C)$$

## Gate Logic: Laws of Boolean Algebra

*Apply the laws and theorems to simplify Boolean equations*

**Example: full adder's carry out function**

$$\text{Cout} = A' B \text{Cin} + A B' \text{Cin} + A B \text{Cin}' + A B \text{Cin}$$

Gate Logic: Laws of Boolean Algebra

*Apply the laws and theorems to simplify Boolean equations*

**Example: full adder's carry out function**

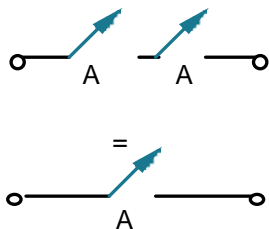
$$\begin{aligned}
 C_{out} &= A' B C_{in} + A B' C_{in} + A B C_{in}' + A B C_{in} \\
 &= A' B C_{in} + A B' C_{in} + A B C_{in}' + \boxed{A B C_{in} + A B C_{in}} \\
 &= A' B C_{in} + \underbrace{A B C_{in} + A B' C_{in} + A B C_{in}' + A B C_{in}} \\
 &= (A' + A) B C_{in} + A B' C_{in} + A B C_{in}' + A B C_{in} \\
 &= (1) B C_{in} + A B' C_{in} + A B C_{in}' + A B C_{in} \\
 &= B C_{in} + A B' C_{in} + A B C_{in}' + \boxed{A B C_{in} + A B C_{in}} \\
 &= B C_{in} + A B' C_{in} + \underbrace{A B C_{in} + A B C_{in}' + A B C_{in}} \\
 &= B C_{in} + A (B' + B) C_{in} + A B C_{in}' + A B C_{in} \\
 &= B C_{in} + A (1) C_{in} + A B C_{in}' + A B C_{in} \\
 &= B C_{in} + A C_{in} + A B (C_{in}' + C_{in}) \\
 &= B C_{in} + A C_{in} + A B (1) \\
 &= B C_{in} + A C_{in} + A B
 \end{aligned}$$

*identity*

*associative*

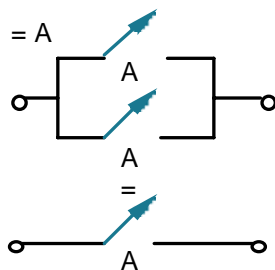
# Gate Logic: Switching Equivalents

$$A \cdot A = A$$

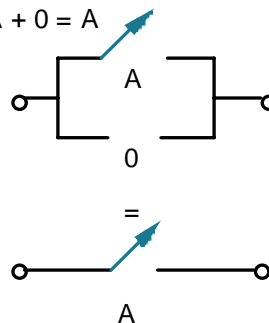


**Idempotent Laws**

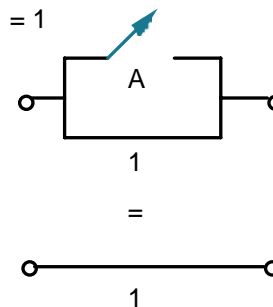
$$A + A = A$$



$$A + 0 = A$$

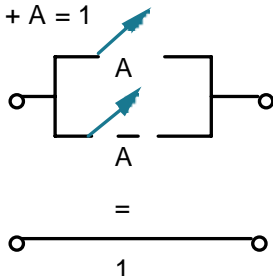


$$A + 1 = 1$$



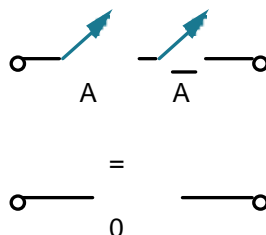
**Identity Laws**

$$A + \bar{A} = 1$$

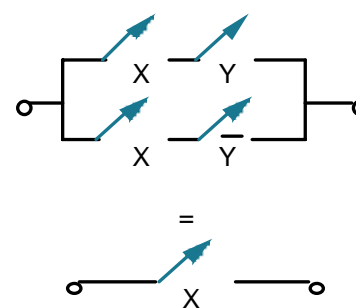


**Complementarity Laws**

$$A \cdot \bar{A} = 0$$

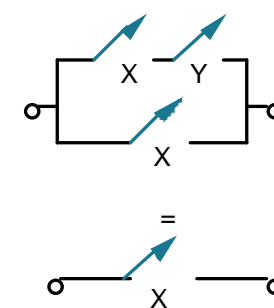


$$XY + X\bar{Y} = X$$



**Simplification Theorems**

$$X + XY = X$$



## Gate Logic: 2-Level Canonical Forms

Truth table is the unique signature of a Boolean function

Many alternative expressions (and gate realizations) may have the same truth table

Canonical form: standard form for a Boolean expression provides a unique algebraic signature

### *Sum of Products Form*

also known as disjunctive normal form, minterm expansion

A	B	C	F	$\bar{F}$
0	0	0	0	1
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	1	0
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

$$F = A' B C + A B' C' + A B' C + A B C' + A B C$$

0 1 1      1 0 0      1 0 1      1 1 0      1 1 1

$$F' = A' B' C' + A' B' C + A' B C'$$